

On Existence of 2-partition of a Tree, Which Obeys the Given Priority

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Abstract

A necessary and sufficient condition is obtained for the problem of such partitioning of the set of vertices of a tree G into two disjoint sets V_1 and V_2 , which, for a given function $p : V(G) \rightarrow \{-1, 0, 1\}$ with some special restriction, satisfies the condition $|\lambda(v) \cap V_1| - |\lambda(v) \cap V_2| = p(v) \cdot (|\{v\} \cap V_1| - |\{v\} \cap V_2|)$ for any vertex v of G , where $\lambda(v)$ is the set of all vertices of G adjacent to v .

We consider finite, undirected graphs without loops or multiple edges. Let $V(G)$ and $E(G)$ denote the sets of vertices and edges of a graph G , respectively. If $v \in V(G)$ then $ex_G(v)$ denotes the eccentricity of a vertex v in a graph G . For a graph G let $\Delta(G)$ be the greatest degree of a vertex of G . Let $\rho_G(x, y)$ denote the distance between the vertices $x \in V(G)$ and $y \in V(G)$ in a graph G . For $v \in V(G)$ let's denote $\lambda(v) \equiv \{\omega \in V(G) / (\omega, v) \in E(G)\}$.

A function $f : V(G) \rightarrow \{0, 1\}$ is called 2-partition of a graph G .

A function $p : V(G) \rightarrow \{-1, 0, 1\}$ is called a priority in a graph G if the following condition holds: for $\forall x \in V(G)$ $p(x) = 0$ iff $d_G(x)$ is even.

We'll say that 2-partition f of a graph G obeys the priority p iff for any vertex $v \in V(G)$

$$|\{\omega \in \lambda(v) / f(\omega) = f(v)\}| - |\{\omega \in \lambda(v) / f(\omega) = 1 - f(v)\}| = p(v).$$

Non-defined concepts can be found in [1, 2, 3, 4, 5].

For any $n \in \{0, 1\}$, an arbitrary function $g : X_g \rightarrow \{0, 1\}$ and any set $X \subseteq X_g$, denote:

$$P(X, g, n) \equiv |\{\nu \in X / g(\nu) = n\}| - |\{\nu \in X / g(\nu) = 1 - n\}|.$$

Obviously, the definition of 2-partition, obeying the given priority, can be rewritten in the following way.

We'll say that 2-partition f of a graph G obeys the priority p iff for any vertex $v \in V(G)$

$$P(\lambda(v), f, f(v)) = p(v).$$

Let $x \in V(G)$ be an arbitrary vertex of a tree G .

We define the subset $N_i(x)$ of the set $V(G)$, where $0 \leq i \leq ex_G(x)$, as follows:

$$N_i(x) \equiv \{z \in V(G) / \rho_G(x, z) = i\}.$$

Obviously, for any $u \in N_i(x)$, where $1 \leq i \leq ex_G(x)$, there exists a single vertex $u^{(-1)} \in N_{i-1}(x)$ satisfying the condition $(u, u^{(-1)}) \in E(G)$.

Let's assume we have some partitioning of the set $V(G) \setminus \{x\}$ into sets $A(x)$, $B(x)$, which satisfies the following conditions:

$$V(G) \setminus \{x\} = A(x) \cup B(x), \quad A(x) \cap B(x) = \emptyset.$$

For $\forall u \in V(G) \setminus N_{ex_G(x)}(x)$ define:

$$\begin{aligned} a(u) &\equiv |N_{p_G(x,u)+1}(x) \cap \lambda(u) \cap A(x)|, \\ b(u) &\equiv |N_{p_G(x,u)+1}(x) \cap \lambda(u) \cap B(x)|. \end{aligned}$$

Note 1. From the definitions of functions a , b it follows that if for $\forall i$, $1 \leq i \leq ex_G(x)$ for all $u \in N_i(x)$ it is already determined whether $u \in A(x)$ or $u \in B(x)$, then for an arbitrary $u \in N_{i-1}(x)$ the values $a(u)$ and $b(u)$ are unambiguously calculated.

Note 2. $a(x) + b(x) = d_G(x)$; for $\forall u \in V(G) \setminus (N_{ex_G(x)}(x) \cup \{x\})$ the equality $a(u) + b(u) + 1 = d_G(u)$ holds.

Now assume that p is a priority in the tree G . Let's inductively define sets $A(x)$ and $B(x)$ as follows:

$$\begin{aligned} N_{ex_G(x)}(x) \cap A(x) &\equiv \{v \in N_{ex_G(x)}(x) / p(v) = 1\}, \\ N_{ex_G(x)}(x) \cap B(x) &\equiv \{v \in N_{ex_G(x)}(x) / p(v) = -1\}. \end{aligned}$$

Let's assume that for i , $2 \leq i \leq ex_G(x)$, the partitioning of $N_i(x)$ is already defined:

$$N_i(x) = (N_i(x) \cap A(x)) \cup (N_i(x) \cap B(x)).$$

It follows from the note 1 that for each $u \in N_{i-1}(x)$ the values of functions a , b can be calculated. Let's define the partitioning of $N_{i-1}(x)$ as follows: for $\forall u \in N_{i-1}(x)$

$$u \in \begin{cases} A(x), & \text{if } a(u) - b(u) - p(u) < 0, \\ B(x), & \text{if } a(u) - b(u) - p(u) > 0. \end{cases}$$

(Note that for any $u \in V(G)$ $a(u) - b(u) - p(u) \neq 0$).

Obviously, under the given definition the following condition is true

$$(N_{i-1}(x) \cap A(x)) \cap (N_{i-1}(x) \cap B(x)) = \emptyset.$$

It is easy to see that the sets $A(x)$ and $B(x)$ are unambiguously defined and, moreover,

$$V(G) \setminus \{x\} = A(x) \cup B(x), \quad A(x) \cap B(x) = \emptyset.$$

Note that we have also defined the following functions

$$a : (V(G) \setminus N_{ex_G(x)}(x)) \rightarrow Z_+, \quad b : (V(G) \setminus N_{ex_G(x)}(x)) \rightarrow Z_+.$$

Further we shall assume, that consideration of any tree G with an arbitrary priority p automatically implies the choice of a vertex $x \in V(G)$, the realization of the partitioning of the set $V(G) \setminus \{x\}$ into sets $A(x)$, $B(x)$ mentioned above and the definition of functions a , b on the set $V(G) \setminus N_{ex_G(x)}(x)$.

Lemma 1. If G is a tree and f - it's 2-partition, obeying the priority p , then for $\forall u \in V(G) \setminus \{x\}$ following properties hold

$$u \in A(x) \Rightarrow f(u^{(-1)}) = f(u), \quad u \in B(x) \Rightarrow f(u^{(-1)}) = 1 - f(u).$$

Proof by the reverse induction on $\rho_G(x, u)$. First of all let's prove the lemma for vertices of the set $N_{\text{ex}_G(x)}(x)$. Let $u \in N_{\text{ex}_G(x)}(x)$ be an arbitrary vertex. Obviously, $d_G(u) = 1$.

Case 1. $u \in A(x)$.

It is clear that $p(u) = 1$. Since f is a 2-partition of a tree G , obeying the priority p and $d_G(u) = 1$, then $f(u^{(-1)}) = f(u)$, which is the statement of the lemma.

Case 2. $u \in B(x)$.

It is clear that $p(u) = -1$. Since f is a 2-partition of the tree G , obeying the priority p , and $d_G(u) = 1$, then $f(u^{(-1)}) = 1 - f(u)$, which is the statement of the lemma. Assume that the lemma holds for all vertices of the set $N_i(x)$, where $2 \leq i \leq \text{ex}_G(x)$. Let's prove the lemma for vertices of the set $N_{i-1}(x)$.

Let $u \in N_{i-1}(x)$ be an arbitrary vertex.

From the inductive assumption it follows that

$$P(\lambda(u) \cap N_i(x), f, f(u)) = a(u) - b(u).$$

Consequently, since f is a 2-partition of the tree G , obeying the priority p , then we obtain

$$\begin{aligned} p(u) &= P(\lambda(u), f, f(u)) = P(\lambda(u) \cap N_i(x), f, f(u)) + \\ &+ P(\lambda(u) \setminus N_i(x), f, f(u)) = a(u) - b(u) + P(\lambda(u) \setminus N_i(x), f, f(u)). \end{aligned} \quad (1)$$

Case 1_{ind}. $u \in A(x)$.

In this case $a(u) - b(u) - p(u) < 0$. Consequently, taking into account the equality (1), we obtain

$$a(u) - b(u) + P(\lambda(u) \setminus N_i(x), f, f(u)) = p(u) > a(u) - b(u).$$

As a result, $P(\lambda(u) \setminus N_i(x), f, f(u)) > 0$. The obtained inequality, taking into account that $\lambda(u) \setminus N_i(x) = \{u^{(-1)}\}$ and $|P(\lambda(u) \setminus N_i(x), f, f(u))| = 1$, implies $f(u^{(-1)}) = f(u)$, which is the statement of the lemma.

Case 2_{ind}. $u \in B(x)$.

In this case $a(u) - b(u) - p(u) > 0$.

Consequently, taking into account the equality (1), we obtain

$$a(u) - b(u) + P(\lambda(u) \setminus N_i(x), f, f(u)) = p(u) < a(u) - b(u).$$

As a result, $P(\lambda(u) \setminus N_i(x), f, f(u)) < 0$. The obtained inequality, taking into account that $\lambda(u) \setminus N_i(x) = \{u^{(-1)}\}$ and $|P(\lambda(u) \setminus N_i(x), f, f(u))| = 1$, implies $f(u^{(-1)}) = 1 - f(u)$, which is the statement of the lemma.

Lemma is proved.

The proof of the lemma implies the following

Corollary 1. If G is a tree and f - it's 2-partition, obeying the priority p , then for any vertex $u \in V(G) \setminus N_{\text{ex}_G(x)}(x)$ the following equality holds

$$p(u) = a(u) - b(u) + P(\lambda(u) \setminus N_{\rho_G(x, u)+1}(x), f, f(u)).$$

Theorem 1. For a given tree G there exists a 2-partition, obeying the priority p , iff for $\forall u \in V(G) \setminus N_{\rho_G(x)}(x)$

$$|a(u) - b(u) - p(u)| = d_G(u) - a(u) - b(u).$$

Proof.

Necessity. Suppose f is a 2-partition, obeying the priority p .

Let $u \in V(G) \setminus N_{\rho_G(x)}(x)$ be an arbitrary vertex. It follows from the corollary 1 and the note 2 that

$$\begin{aligned} |a(u) - b(u) - p(u)| &= \\ &= |a(u) - b(u) - (a(u) - b(u) + P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(x), f, f(u)))| = \\ &= |P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(x), f, f(u))| = d_G(u) - a(u) - b(u). \end{aligned}$$

Sufficiency. Suppose that for any $u \in V(G) \setminus N_{\rho_G(x)}(x)$ $|a(u) - b(u) - p(u)| = d_G(u) - a(u) - b(u)$.

Let's inductively define a function $f: V(G) \rightarrow \{0, 1\}$. Let's set $f(x) \equiv 1$. Let's assume that for all vertices of the set $N_i(x)$, where $0 \leq i \leq \rho_G(x) - 1$, the function f is already defined. Let's define the function f for vertices of the set $N_{i+1}(x)$. For each vertex $u \in N_i(x)$ let's define the function f for vertices of the set $N_{i+1}(x) \cap \lambda(u)$. First of all let's define the function f on vertices of the set $N_{i+1}(x) \cap \lambda(u) \cap A(x)$ by the following way: for $\forall z \in N_{i+1}(x) \cap \lambda(u) \cap A(x)$ set $f(z) \equiv f(u)$.

Now let's define the function f on vertices of the set $N_{i+1}(x) \cap \lambda(u) \cap B(x)$ by the following way: for $\forall z \in N_{i+1}(x) \cap \lambda(u) \cap B(x)$ set $f(z) \equiv 1 - f(u)$. So we have defined the function f on all vertices of the set $N_{i+1}(x)$. Therefore, the function f is defined on whole $V(G)$. Let's check that 2-partition f of the tree G , defined above, obeys the priority p , indeed.

Let $u \in V(G)$ be an arbitrary vertex.

Case 1. $u = x$.

The condition of the theorem and the note 2 imply $|a(x) - b(x) - p(x)| = 0$.

Consequently, from the equality $\lambda(x) = \lambda(x) \cap N_1(x)$ and from the definition of f we obtain

$$P(\lambda(x), f, f(x)) = P(\lambda(x) \cap N_1(x), f, f(x)) = a(x) - b(x) = p(x).$$

Case 2. $u \neq x$. The condition of the theorem and the note 2 imply $|a(u) - b(u) - p(u)| = 1$. Consequently, from the definition of f we obtain

$$\begin{aligned} P(\lambda(u), f, f(u)) &= P(\lambda(u) \cap N_{\rho_G(x,u)+1}(u), f, f(u)) + \\ &+ P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = \\ &= a(u) - b(u) + P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)). \end{aligned} \quad (2)$$

Case 2a. $u \in A(x)$. In this case from the definitions of $A(x)$ and f it follows that $a(u) - b(u) - p(u) < 0$ and $P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = 1$, respectively. Consequently, taking into account $|a(u) - b(u) - p(u)| = 1$, we obtain $a(u) - b(u) - p(u) = -1$. The obtained equality, taking into account the equality (2) and the equality $P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = 1$, implies

$$\begin{aligned} P(\lambda(u), f, f(u)) &= a(u) - b(u) + P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = \\ &= a(u) - b(u) + 1 = p(u). \end{aligned}$$

Case 2b. $u \in B(x)$.

In this case from the definitions of $B(x)$ and f it follows that $a(u) - b(u) - p(u) > 0$ and $P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = -1$, respectively. Consequently, taking into account $|a(u) - b(u) - p(u)| = 1$, we obtain $a(u) - b(u) - p(u) = 1$. The obtained equality, taking into account the equality (2) and the equality $P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = -1$, implies

$$\begin{aligned} P(\lambda(u), f, f(u)) &= a(u) - b(u) + P(\lambda(u) \setminus N_{\rho_G(x,u)+1}(u), f, f(u)) = \\ &= a(u) - b(u) - 1 = p(u). \end{aligned}$$

Theorem is proved.

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Ծառի այնպիսի 2-տրոհման գոյության մասին,
որը ենթարկվում է տրված նախապատվությանը

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Ամսփոփում

Ստացված է անհրաժեշտ և բավարար պայման G ծառի զագաթների բազմության V_1 և V_2 չհատվող ենթաբազմությունների այնպիսի տրոհման գոյությունը պարզելու համար, որ տրված հատուկ սահմանափակումներով ֆունկցիայի համար բավարարվի հետևյալ պայմանը $\nu: V(G) \rightarrow \{-1, 0, 1\}$ ծառի յուրաքանչյուր ν զագաթի համար $|\lambda(v) \cap V_1| - |\lambda(v) \cap V_2| = \nu(v) \cdot (|\{v\} \cap V_1| - |\{v\} \cap V_2|)$, որտեղ $\lambda(v)$ -ով նշանակված է v -ին կից զագաթների բազմությունը: